Using Regression Techniques to Predict Large Data Transfers

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Abstract

The recent proliferation of Data Grids and the increasingly common practice of using

resources as distributed data stores provide a convenient environment for communities

of researchers to share, replicate, and manage access to copies of large datasets. This

has led to the question of which replica can be accessed most efficiently. In such

environments, fetching data from one of the several replica locations requires accurate

predictions of end-to-end transfer times. The answer to this question can depend on

many factors, including physical characteristics of the resources and the load behavior

on the CPUs, networks, and storage devices that are part of the end-to-end data path

linking possible sources and sinks.

Our approach combines end-to-end application throughput observations with network

and disk load variations and captures whole-system performance and variations in load

patterns. Our predictions characterize the effect of load variations of several shared

devices (network and disk) on file transfer times. We develop a suite of univariate and

multivariate predictors that can use multiple data sources to improve the accuracy of

the predictions as well as address Data Grid variations (availability of data and

sporadic nature of transfers). We ran a large set of data transfer experiments using

GridFTP and observed performance predictions within 15% error for our testbed sites,

which is quite promising for a pragmatic system.

Keywords: *Grids, data transfer prediction, replica selection.*

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1. Introduction

As the coordinated use of distributed resources, or Grid computing, becomes more commonplace, basic resource usage is changing. Many recent applications use Grid systems as distributed data stores [DataGrid02, GriPhyN02, HSS00, LIGO02, MMR+01, NM02], where pieces of large datasets are replicated over several sites. For example, several high-energy physics experiments have agreed on a tiered Data Grid architecture [HJS+00, Holtman00] in which all data (approximately 20 petabytes by 2006) is located at a single Tier 0 site; various (overlapping) subsets of this data are located at national Tier 1 sites, each with roughly one-tenth the capacity; smaller subsets are cached at smaller Tier 2 regional sites; and so on. Therefore, any particular dataset is likely to have replicas located at multiple sites.

Different sites may have varying performance characteristics because of diverse storage system architectures, network connectivity features, or load characteristics. Users (or brokers acting on their behalf) may want to be able to determine the site from which particular data sets can be retrieved most efficiently, especially as data sets of interest tend to be large (1–1000 MB). It is this *replica selection* problem that we address in this paper.

Since large file transfers can be costly, there is a significant benefit in selecting the most appropriate replica for a given set of constraints [ACF+02, VTF01]. One way a more intelligent replica selection can be achieved is by having replica locations expose performance information about past data transfers. This information can, in theory, provide a reasonable approximation of the end-to-end throughput for a particular transfer. It can then be used to make predictions about the future behavior between the sites involved. In our work we use GridFTP [AFN+01], part of the Globus ToolkitTM [FK98, Globus02] for moving data. We are interested in predicting the performance of GridFTP transfers for large files across the Grid.

In this paper we present two- and three-datastream predictions using regression techniques. We start by deriving predictions from past history of GridFTP transfers in isolation. We build a suite of univariate predictors comprising simple mathematical models such as mean- and median-based tools that are easy to implement and achieve acceptable levels of accuracy. We then present a detailed analysis of several variations of our univariate forecasting tools and information on GridFTP logs.

The univariate models do not achieve better prediction accuracy because they fail to account for the sporadic nature of data transfers in Grid environments. Hence, predictions based on simple log data may not contain recent enough information on current system trends. We need to be able to derive forecasts from several combinations of "currently available" data sources.

To address this need, we use both log data and periodic data to explain the behavior of key components in the end-to-end data path. We use the additional datastreams of network and disk behavior to illustrate how additional data can be exploited in predicting the behavior of large transfers. We present an in-depth study of these data sources and our multivariate forecasting tools, including information about data formats, lifetime, time/space constraints, correlation, statistical background on our regression tools, and the pros and cons of this approach.

We then evaluate our prediction approaches using several different metrics. Comparing the normalized percentage errors of our various predictions, we find that the univariate predictions have error rates of at most 25% and that all the univariate predictors performed similarly. With multivariate predictions, we observed that combining GridFTP logs and disk throughput observations provided us with gains of up to 4% when compared with the best of univariate predictors. Combining logs with network throughput data provides further gains up to 6%, and predictions based on all three data sources had up to 9% reduction in error. To study the degree of variance in error rates, we computed confidence levels and observed that the variance is smaller with more accurate predictors for the sites we examined. We further developed a triplet metric

comprising the throughput, percentage error rate, and confidence level as a measure of a given site's predictive merit.

2. Related Work

The goal of this work is to obtain accurate predictions of file transfer times between a storage system and a client. Achieving this can be challenging because numerous devices are involved in the end-to-end path between the source and the client, and the performance of each (shared) device along the end-to-end path may vary in unpredictable ways.

One approach to predicting this information is to construct performance models for each system component (CPUs at the level of cache hits and disk access, networks at the level of the individual routers, etc.) and then to use these models to determine a schedule for all data transfers [SC00], similar to classical scheduling [Adve93, Cole89, CQ93, Crovella99, ML90, Schopf97, TB86, ZLP96]. In practice, however, it is often unclear how to combine this data to achieve accurate end-to-end measurements. Also, since system components are shared, their behavior can vary in unpredictable ways [SB98]. Further, modeling individual components in a system may not capture the significant effects these components have on each other, thereby leading to inaccuracies [GT99].

Alternatively, observations from past application performance of the entire system can be used to predict end-to-end behavior. The use of whole-system observation has relevant properties for our purposes. These predictions can, in principle, capture both evolution in system configuration and temporal patterns in load. A by-product of capturing entire system evolution is enhanced transparency, in that we can construct such predictions without detailed knowledge of the underlying physical devices. This technique is used by Downey [Downey97] and Smith et al. [SFT98] to predict queue wait times and by numerous tools (Network Weather Service [Wolski98], NetLogger

[NetLogger02], Web100 [Web100Project02], iperf [TF01], and Netperf [Jones02]) to predict the network behavior of small file transfers.

Although tools such as the Network Weather Service (NWS) measure and predict network bandwidth, a substantial difference in performance can arise between a small NWS probe (lightweight with 64 KB size) and an actual file transfer using GridFTP (with tuned TCP buffers and parallelism). We show this in Figure 1, which depicts 64 KB NWS measurements, indicating that the bandwidth is about 0.3 MB/sec, and end-to-end GridFTP measurements, indicating a significantly higher transfer rate. In this case, NWS by itself is not sufficient to predict end-to-end GridFTP throughput. In addition, we see a much larger variability in GridFTP measurements, ranging from 1.5 to 10.2 MB/sec (because of different transfer sizes and also load variations in the end-to-end components), so that it is unlikely that a simple data transformation will improve the resulting prediction.

The univariate predictors presented in this work are similar to the basic predictors used by NWS and similar tools to predict the behavior of time series data. Because our data traces are not periodic in nature, however, we also use predictions based on multiple datastreams. This approach is similar to work done by Faerman et al. [FSW+99], which used the NWS and adaptive linear regression models for the Storage Resource Broker [BMR+98] and SARA [SARA02]. Faerman and his colleagues compared transfer times obtained from a raw bandwidth model (Transfer-Time = ApplicationDataSize/NWS-Probe-Bandwidth, with 64 KB NWS probes) with predictions from regression models and observed accuracy improvements ranging from 20% to almost 100% for the sites examined. The work presented here goes beyond that work, however, by exploring several filling techniques to mitigate adverse effects of sporadic transfers.

Swany and Wolski have also approached multivariate predictors by constructing cumulative distribution functions of past history and deriving predictions from them as an alternative to regressive models. This approach has been demonstrated for 16 MB HTTP transfers with improved prediction accuracy when compared with their

univariate prediction approach [SW02]. Further, they have applied their models to our datasets, comprising various file sizes, and have observed comparable prediction accuracy.

3. Data Sources

In this section, we describe our three primary data sources. We use the GridFTP server to perform our data transfers and log its behavior every time a transfer is made, thereby recording the end-to-end transfer behavior. Since these events are very sporadic in nature, however, we also need to capture data about the current environment to have accurate predictions. Hence, we use the Network Weather Service probe data as an estimate of bandwidth for small data transfers and the iostat disk throughput data to measure disk behavior.

3.1. GridFTP Logs

GridFTP [AFN+01] is part of the Globus ToolkitTM [FK98, Globus02] and is widely used as a secure, high-performance data transfer protocol [ACF+02, AFN+01, DataGrid02, GriPhyN02]. It extends standard FTP implementations with several features needed in Grid environments, such as security, parallel transfers, partial file transfers, and third party transfers.. We instrumented the GT 2.0 wuftp-based GridFTP server to log the source address, file name, file size, number of parallel streams, stripes, TCP buffer size for the transfer, start and end timestamps, nature of the operation (read/write), and logical volume to/from which file was transferred (Table 1) [VSF02].

The GridFTP monitoring code is nonintrusive. The majority of the overhead is in the timing routines, with a smaller percentage spent gathering the information mentioned above and performing a write operation. The entire logging process consumes on average of approximately 25 milliseconds per transfer, which is insignificant compared with the total transfer time. The NWS and disk throughput logging consume much less time (only a few milliseconds) when compared with GridFTP logging.

Although each log entry (from all three datastreams) is well under 512 bytes, transfer logs can grow quickly in size at a busy site. We do not currently implement a log management scheme, but it would be straightforward to use a circular buffer, such as in the NWS. An alternative strategy used by NetLogger is to flush the logs to persistent storage (either disk or network) and restart logging.

3.2. **NWS**

The Network Weather Service [Wolski98] monitors the behavior of various resource components by sending out lightweight probes or querying system files at regular intervals. NWS sensors exist for components such as CPU, disk, and network. We used the network bandwidth sensor with 64 KB probes to estimate the current network throughput. NWS throughput measurements, although not representative of the transfer bandwidth obtainable for large files (10 MB to 1 GB), are representative of the network link characteristics. Further, NWS is intended to be a lightweight, noninvasive monitoring system whose measurements can then be extrapolated to specific cases such as ours.

3.3. Iostat

Traditionally, in large wide-area transfers, network transport has been considered to weigh heavily on the end-to-end throughput achieved. Current trends in disk storage and networking, however, suggest that disk accesses will factor rather strongly in the future. Network throughput is far outpacing advances in disk speeds. Therefore, as link speeds increase, the network latency significantly drops, and disk accesses are likely to become the bottleneck in large file transfers across the Grid [GS00].

To address this issue, we include disk throughput data in our prediction approach. The iostat tool is part of the sysstat [SYSSTAT02] system-monitoring suite and collects disk I/O throughput data by monitoring the blocks read/written from/to a particular

disk. Iostat can be configured to periodically monitor disk transfer rates, block read/write rates, and so forth of all physically connected disks. We use the disk transfer rate that represents the throughput of the disk.

3.4. Correlation

A key step in analyzing whether a combination of datastreams will result in better predictions is to evaluate how highly correlated they are. The correlation coefficient is a measure of the linear relationship between two variables and can have a value between -1.0 and +1.0 depending on the strength of the relation. A coefficient near zero suggests that the variables may not be linearly related, although they may exhibit nonlinear dependencies [Edwards84, OM88]. The correlation coefficient for two datastreams G and N is computed by using the formula

$$corr = \frac{\sum NG - (\sum N\sum G/size)}{\sqrt{(\sum G^2 - (\sum G)^2/size)} \sqrt{(\sum N^2 - (\sum N)^2/size)}}$$

where "size" is the number of values in the data stream.

We compute the rank-order correlation for each of our datasets. Rank correlation provides a distribution-free, nonparametric alternative to determine whether the observed correlation is significant [Edwards84]. Rank correlation converts data to ranks by assigning a specific rank to each value in the datastream, as determined by the position of the value when the datastream is sorted. Table 2 shows a tabulated listing of the 95% confidence interval for the correlation coefficients for the three datasets we collected between our transfer points. The confidence interval denotes that the correlation for 95% of the sample falls within a certain upper and lower limit. We can see a moderate correlation between GridFTP, NWS, and disk throughput datastreams.

4. Predictors

We evaluated a wide set of prediction techniques for wide-area data transfers. This section presents the univariate predictions and the multivariate prediction techniques we used in our experiments.

4.1. Univariate Predictors

In this section we describe some of the predictors we developed, categorize possible approaches by basic mathematical techniques, and detail the pros and cons of each technique.

4.1.1. Mathematical Functions

Mathematical functions for predictions are generally grouped into mean-based, medianbased, and autoregressive techniques. We use several variations of each of these models in our experiments.

Mean-based, or averaging, techniques are a standard class of predictors that use arithmetic averaging (as an estimate of the mean value) over some portion of the measurement history to estimate future behavior. The general formula for these techniques is the sum of the previous n values over the number of measurements. Mean-based predictors are easy to implement and impose minimally on system resources.

A second class of standard predictors is based on evaluating the median of a set of values. Given an ordered list of t values, if t is odd, the median is the (t+1)/2 value; if t is even, the median is half of the t/2 value added with the (t+1)/2 value. Median-based predictors are particularly useful if the measurements contain randomly occurring asymmetric outliers that are rejected. However, they lack some of the smoothing that occurs with a mean-based method, possibly resulting in forecasts with a considerable amount of jitter [HP91].

The third class of common predictors is autoregressive models [GP94, HP91, Wolski98]. We use an autoregressive integrated moving average (ARIMA) model technique that is constructed using the equation

$$G^{|} = a + bG_{t-1}$$

where G^{l} is the GridFTP prediction for time, t, G_{t-1} is the previous data occurrence, and a and b are the regression coefficients that are computed based on past occurrences of G using the method of least squares. This approach is most appropriate when there are at least fifty measurements and the data is measured with equally spaced time intervals. Our data does not meet these constraints, but we include this technique to do a full comparison. The main advantage of using an ARIMA model is that it gives a weighted average of the past values of the series, thereby possibly giving a more accurate prediction. However, in addition to requiring a larger data set than the other techniques to achieve a statistically significant result, the model can have a much greater computational cost.

4.1.2. Context-Insensitive Variants

More recent values are often better predictors of future behavior than an entire dataset, no matter which mathematical technique is used to calculate a prediction. Hence, many different variants exist in selecting a set of recent measurements to use in a prediction calculation.

The fixed-length, or sliding window, average is calculated by using only a set number of previous measurements to calculate the average. The number of measurements can be chosen statically or dynamically depending on the system. We use only static selection techniques in this work. Options for dynamically selecting window size are discussed in [Wolski98]. The degenerative case of this strategy involves using only the last measurement to predict the future behavior. Work by Downey and Harchol-Balter [HD96] shows that this is a useful predictor for CPU resources, for example.

Instead of selecting the number of recent measurements to use in a prediction, we also consider using only a set of measurements from a previous window of time. Unlike other systems where measurements are taken at regular intervals [DO00, Wolski98], our measurements can be spaced irregularly. Using temporal windows for irregular samples can reflect trends more accurately than selecting a specific number of previous measurements because they capture recent fluctuations, thereby helping to ensure that recent (and, one hopes, more predictive) data is used. Much as the number of measurements included in a prediction can be selected dynamically, the window of time used can be decided dynamically.

As shown in Table 3, we use fixed-length sets of the last 1 (last value), 5, 15, and 25 measurements. We use temporal-window sets of data of the last 5 hours, 15 hours, 25 hours, 5 days, and 10 days. We consider both mean-based and median-based predictors over previous *n* measurements; mean-based predictors over the previous 5, 15, and 25 hours; and autoregression (AR) over the previous 5 and 10 days, since this function requires a much larger dataset to produce accurate predictions than our other techniques.

4.1.3. Context-Sensitive Variants

Filtering a data set to eliminate unrelated values often results in a more accurate prediction. For example, a prediction of salary is more accurate when factors such as previous training, education, and years at the position are used to limit the dataset of interest.

With the GridFTP monitoring data, initial results showed that file transfer rates had a strong correlation with file size. Studies of Internet traffic have also revealed that small files achieve low bandwidths whereas larger files tend to have high bandwidths [BMK96, CSA98, GM01]. This difference is thought to be primarily due to the startup overhead associated with the TCP start mechanism that probes the bandwidth at connection startup. Recent work has focused on class-based isolation of TCP flows

[YM01] and on startup optimizations [ZQK00, ZQK99] to mitigate this problem. As a proof of concept, we found 5%–10% improvement on average when using file-size classification instead of the entire history file to calculate a prediction. This is discussed in Section 5.

For our GridFTP transfer data we ran a series of tests between our testbed sites to categorize the data sizes into a small number of classes. We categorized our data into four sets: 0–50 MB, 50–250 MB, 250–750 MB, and more than 750 MB based on the achievable bandwidth. We note that these classes apply to the set of hosts for our testbed only; further work is needed to generalize this notion.

4.2. Multivariate Predictors

The obvious downside of univariate predictors has nothing to do with the predictors themselves but more so with the nature of data transfers on the Grid. Because of the sporadic nature of transfers, predictors based on log data alone may fail to factor in current system trends and fluctuations. To mitigate the adverse effects of this problem, we introduce other "steady" and periodic datastreams to describe the behavior of components in the end-to-end data path, capturing current trends.

We developed a set of multivariate predictors using regression models to predict from a combination of several data sources – GridFTP log data and network load data, GridFTP log data and disk load data, or a combination of all three. The datastreams require some preprocessing before the regression techniques can be applied to them. This includes time matching the data streams and filling-in techniques.

4.2.1 Matching

Our three data sources (GridFTP, disk I/O, and NWS network data) are collected exclusive of each other and rarely have the same timestamps. To use regressive models on the data streams, however, we need to have a one-to-one mapping for the values in

each stream. Hence, we are required to match values from the three sets such that for each GridFTP value, we find disk I/O and network observations that were made around the same time.

For each GridFTP data point (T_g, G) , we match a corresponding disk load (T_D, D) and NWS data point (T_N, N) such that T_N and T_D are the closest to T_G . By doing this, the triplet (N_i, D_j, G_k) represents an observed end-to-end GridFTP throughput (G_k) resulting from a data transfer that occurred with the disk load (D_j) and network probe value (N_i) .

At the end of the matching process, the three datastreams have been combined into the sequence that looks like

$$(N_i, D_j, G_k)(N_{i+1}, D_{j+1},)...(N_{i+m}, D_{j+m}, G_{k+1}),$$

where G_k , and G_{k+1} are two successive GridFTP file transfers, N_i and N_{i+m} are NWS measurements, and D_j and D_{j+m} are disk load values that occurred in the same timeframe as the two GridFTP transfers. The sequence also consists of a number of disk load and NWS measurements between the two transfers for which there are no equivalent GridFTP values, such as $(N_{i+1}, D_{j+1}, _)$. Note that these interspersed network and disk load values are time-aligned. Also note that we have described the matching process with reference to all three data sources. In the case where a prediction uses a different number of datastreams, similar matching techniques can be employed.

4.2.2. Filling-in Techniques

After matching the datastreams, we need to address the tuples that do not have values for the GridFTP data – that is, the NWS data or disk I/O data collected in between the sporadic GridFTP transfers. Regression models expect a one-to-one mapping between the data values, so we can either discard the network and I/O data for which there are no equivalent GridFTP data (our NoFill technique, Figure 2) or fill in synthetic transfer values using either an average over the past day's data (Avg), or the last value (LV). Once filled in, the sequence is as follows:

$$(N_i, D_j, G_k)(N_{i+1}, D_{j+1}, G_{Fill})...(N_{i+m}, D_{j+m}, G_{k+1})$$

where G_{Fill} is the synthetic GridFTP value. Data, once matched and filled in, is fed to regression models (Figure 3).

4.2.3. Linear Regression

Linear regression attempts to build linear models between dependent and independent variables. The following equation builds linear models between several independent variables N_1 , N_2 , ..., N_k and dependent variable G as follows:

$$G^{\dagger} = a + b_1 N_1 + b_2 N_2 + ... + b_k N_k$$

where G^{\dagger} is the prediction of the observed value of G for the corresponding values of N_1 , N_2 , ..., N_k . The coefficients a, b_1 , b_2 , and b_k are calculated by using the method of least squares [Edwards84]. For our case, we built linear models between NWS (N), disk (D), and GridFTP (G) data as explained above, with N and D as independent variables.

4.2.4. Polynomial Regression Models

To improve prediction accuracy, we also developed a set of nonlinear models adding polynomial terms to the linear equation. For instance, a quadratic model is as follows:

$$G^{\dagger}=a+b_1N+b_2N^2$$
.

Cubic and quartic models have additional terms b₃N³ and b₄N⁴, respectively. Similar to the linear model, the coefficients in quadratic, cubic, and quartic models b₂, b₃, and b₄ are computed by using the method of least squares. Adding polynomial terms to the regression model can reach a saturation point (no significant improvement in prediction accuracy observed), suggesting that a particular model sufficiently captures the relationship between the two variables [OM88, Pankratz91]. Figure 4 shows a bar graph that compares error, complexity of algorithm, and components included for the site pair, Lawrence Berkeley and Argonne National Laboratories.

5. Measurements and Evaluation

We evaluated the performance of our regression techniques on datasets collected over three distinct two-week durations: August 2001, December 2001, and January 2002. In the following subsections we describe the experimental setup, prediction error calculations, and the results obtained from these datasets.

5.1. Experimental Setup

The experiments we ran consisted of controlled GridFTP transfers, NWS network sensor measurements, and disk throughput monitoring between four sites in our testbed (Figure 5): Argonne National Laboratory (ANL), the University of Southern California Information Sciences Institute (ISI), Lawrence Berkeley National Laboratory (LBL) and the University of Florida at Gainesville (UFL).

GridFTP experiments included transfers comprising several file sizes ranging from 10 MB to 1 GB, performed at random time intervals within 12-hour periods. We calculated buffer sizes by using the formula

RTT * "bottleneck bandwidth in the link"

with roundtrip times (RTT) values obtained from ping and with bottleneck bandwidth obtained by using iperf [TF01]. Figure 5 shows the roundtrip times and bottleneck bandwidth for our site pairs. Our GridFTP experiments were performed with tuned TCP buffer settings (1 MB based on the bandwidth delay product) and eight parallel streams to achieve enhanced throughput. Logs of these transfers were maintained at the respective sites and can be found at [Traces02].

Configuring NWS among a set of resources involved setting up a nameserver and memory to which sensors at various sites registered and logged measurements [Wolski98]. In our experiments, we used ANL as a registration and memory resource. NWS network monitoring sensors between these sites were set up to measure bandwidth every five minutes with 64 KB probes.

Disk I/O throughput data was collected by using the iostat tool logging transfer rates every five minutes. Logs were maintained at the respective servers.

For each data set and predictor, we used a 15-value training set; that is, we assumed that at the start of a predictive technique there were at least 15 GridFTP values in the log file (approximately two days worth of data).

5.2. Metrics

We calculate the prediction accuracy using the normalized percentage error calculation:

% Error =
$$\frac{\sum |Measured_{BW} - Predicted_{BW}|}{(size * Mean_{BW})} * 100,$$

where "size" is the total number of predictions and the Mean is the average measured GridFTP throughput. In this subsection we show results based on the August 2001 dataset. Complete results for all three datasets can be found in the appendix and at [Traces02].

In addition to evaluating the error of our predictions, we evaluate information about the variance in the error. Depending on the use case, a user may be more interested in selecting a site that has reasonable performance bandwidth estimates with a relatively low prediction error than in selecting a resource with higher performance estimates and a possibly much higher error in prediction. In such cases, it can be useful if the forecasting error can be stated with some confidence and with a maximum/minimum variation range. These limits can also, in theory, be used as catalysts for corrective measures in case of performance degradation.

In our case, we can also use these limits to verify the inherent cost of accuracy of the predictors. By comparing the confidence intervals of these prediction error rates, we can determine whether the accuracy achieved is at the cost of greater variability, in which case there is little gain in increasing the component complexity of our prediction approach.

Thus, for any predictor (for any site pair and a given dataset), the information denoted by the following triplet can be used as a metric to gauge its accuracy:

Accuracy-Metric = [PredictedThroughput, AvgPast % Error-Rate, ConfidenceLimit], where PredictedThroughput is the predicted GridFTP value (higher the better), with a certain percentage prediction error (the lower the better) and a percentage confidence interval for the error (the smaller the better).

5.3. Univariate Predictor Performance

Figure 6 shows bar charts of percentage error for our various univariate predictors at the various site pairs. The major result from these predictions is that even simple techniques are "at worst" off by about 25%, quite respectable for pragmatic prediction systems.

Figure 7 shows the result of sorting the data by file size, since GridFTP throughput varied with transfer file sizes. We grouped several file sizes into categories: 0–50 MB as 10M, 50–250 MB as 100M, 250–750 MB as 500M, and more than 750 MB as 1G, based on the achievable bandwidth. We observe almost up to 10% increase in accuracy with filtering context sensitive filtering.

Figure 8 shows the relative performance of the predictors to determine which predictor performed better by computing the best and worst predictor for each data transfer. On average, predictors that had a high "best" percentage also had a high "worst" percentage.

In general, for our univariate predictors,, we did not see a noticeable advantage of limiting either average or median techniques using a sliding window or time frames. The ARIMA models did not see improved performance for our data, although they are significantly more expensive. This is likely due to the irregular nature of our data.

5.4. Multivariate Predictor Performance

Table 4 shows the performance gains of using a regression prediction with GridFTP and NWS network data (G+N) over using the GridFTP log data univariate predictor in isolation (first two shaded columns in the table). We use the moving average (AVG25) as a representative of univariate predictor performance. For our datasets, we observed a 4% to 6% improvement in prediction accuracy when the regression techniques with LV or AVG filling were used. Regression with NoFill (throwing away the unmatched GridFTP data) shows no significant improvement when compared with univariate predictors.

Table 4 also shows that including disk I/O component load variations in the regression model provides us with gains of 2% to 4% (G+D Avg) when compared with AVG25 (first and third shaded columns in the table). Different filling techniques (G+D Avg and G+D LV) perform similarly, and again NoFill shows no improvement, or even a decrease in accuracy, when compared with univariate predictors.

Comparing the second and third block of data in Table 4 shows that all variations of predictors using NWS data (G+N) perform better than predictors using disk I/O data (G+D) in general. This observation agrees with our initial measurements that only 15% to 30% of the total transfer time is spent in I/O, while the majority of the transfer time (in our experiments) is spent performing network transport.

When we include both disk I/O and NWS network data in the regression model (G+N+D) along with GridFTP transfer logs, we see prediction error drop to 8% to 17% and up to 3% improvement when compared with G+N (second and fourth shaded columns in Table 4) and a 6% improvement over G+D (third and fourth shaded columns in Table 4). Overall, we see up to 9% improvement when we compare G+N+D with the original univariate prediction based on AVG25.

Figure 9a compares the average prediction error for Moving Avg, G+D Avg, G+N Avg, and G+N+D Avg for all of our site pairs (represents the shaded columns in Table 4) and also presents 95% confidence limits for our prediction error rates. The prediction accuracy trend is as follows:

Moving
$$Avg < (G+D Avg) < (G+N Avg) < (G+N+D Avg)$$

Figure 9b shows that the confidence interval (the variance in the error) does in fact reduce with more accurate predictors, but the reduction is not significant for our datasets.

Figure 10 depicts the performance of predictors G+D Avg, G+N Avg and G+N+D Avg. The predictors closely track the measured GridFTP values. Predictions were obtained by using regression equations that were computed for each observed network or disk throughput value.

For our datasets, we observed no noticeable improvement in prediction accuracy by using polynomial models for our site pairs. We studied the effects of polynomial regression on all our multivariate tools (G+D, G+N and G+N+D). Figure 11 shows the performance of linear, quadratic, cubic, and quartic regression models for various site pairs for the G+D Avg predictor. All our models performed similarly. On average, regression-based predictors with filling took approximately 10 msec for a dataset size of 50 GridFTP and 1500 NWS values.

6. Conclusions

In this paper we describe the need for predicting the performance of GridFTP data transfers in the context of replica selection in Data Grids. We show how bulk data transfer predictions can be derived and how its accuracy can be improved by including information on current system/network trends. Further, we argue how data transfer predictions can be constructed using several combinations of datasets. We detail the development of a suite of univariate and multivariate predictors that satisfy the specific constraints of Data Grid environments. We examine a series of simple univariate predictors that are lightweight and use means, medians, and autoregressive techniques. We also use more complex regression analysis for multivariate predictors. To mitigate the adverse effects of sporadic transfers, multivariate predictors use several filling-in techniques such as last value (LV) and average (AVG). We observe that multivariate

predictors with filling offer considerable benefits (up to 9%) when compared with univariate predictors.

Appendix

Tables 5 and 6 show the performance gains of using a regression prediction with GridFTP and NWS network data (G+N) over using the GridFTP log data univariate predictor in isolation for the December 2001 and January 2002 datasets. Behaviors of both univariate and multivariate predictors are similar to those exhibited in the August 2001 dataset (Table 4). In general, we observe performance improvements in using regression-based filling predictors and prediction error reduces with the addition of disk and network data.

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Table 1: Sample set from a log of file transfers between Argonne and Lawrence Berkeley National Laboratories. The bandwidth values logged are sustained measures through the transfer. The end-to-end GridFTP bandwidth is obtained by the formula BW = file size / transfer time.

Source IP	File Name	File Size	Volume	StartTime	EndTime	TotalTime	Bandwidth	Read/Write	Streams	TCP-
		(Bytes)		(Timestamp)	(Timestamp)	(Seconds)	(KB/Sec)			Buffer
140.221.65.69	/home/ftp/vazhkuda/10 MB	10240000	/home/ftp	998988165	998988169	4	2560	Read	8	1000000
140.221.65.69	/home/ftp/vazhkuda/25 MB	25600000	/home/ftp	998988172	998988176	4	6400	Read	8	1000000
140.221.65.69	/home/ftp/vazhkuda/50 MB	51200000	/home/ftp	998988181	998988190	9	5688	Read	8	1000000
140.221.65.69	/home/ftp/vazhkuda/100 MB	102400000	/home/ftp	998988199	998988221	22	4654	Read	8	1000000
140.221.65.69	/home/ftp/vazhkuda/250 MB	256000000	/home/ftp	998988224	998988256	33	8000	Read	8	1000000
140.221.65.69	/home/ftp/vazhkuda/500 MB	512000000	/home/ftp	998988258	998988335	67	7641	Read	8	1000000
140.221.65.69	/home/ftp/vazhkuda/750 MB	768000000	/home/ftp	998988338	998988425	97	7917	Read	8	1000000
140.221.65.69	/home/ftp/vazhkuda/1 GB	1024000000	/home/ftp	998988428	998988554	126	8126	Read	8	1000000

Table 2: 95% Confidence for the upper and lower limits of the rank-order correlation coefficient for the GridFTP, NWS, and disk I/O datasets between four sites in our testbed. Denotes coefficients for our three datasets.

	GridFTP and NWS							GridFTP and Disk I/O					
	Aug'01		Dec	'01	Jar	'02	Aug'01		Dec'01		Jan'02		
	Upper	Lower	Upper	Lower	Upper	Lower	Upper	Lower	Upper	Lower	Upper	Lower	
LBL-ANL	0.8	0.5	0.5	0.3	0.6	0.2	0.6	0.1	0.5	0.2	0.5	0.1	
LBL-UFL	0.7	0.5	0.7	0.4	0.6	0.1	0.5	0.2	0.5	0.3	0.5	0.3	
ISI-ANL	0.8	0.5	0.6	0.4	0.7	0.3	0.5	0.2	0.6	0.4	0.6	0.3	
ISI-UFL	0.9	0.4	0.6	0.2	0.5	0.1	0.5	0.1	0.6	0.3	0.5	0.2	
ANL-UFL	0.5	0.2	0.6	0.2	0.6	0.1	0.5	0.2	0.4	0.1	0.4	0.2	

Table 3: Context-insensitive predictors used

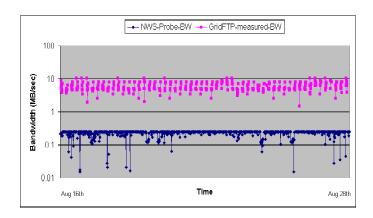
	Average based	Median based	Autoregression
All data	AVG	MED	AR
Last 1 Value	LV		
Last 5 Values	AVG5	MED5	
Last 15 Values	AVG15	MED15	
Last 25 Values	AVG25	MED25	
Last 5 Hours	AVG5hr		
Last 15 Hours	AVG15hr		
Last 25 Hours	AVG25hr		
Last 5 Days			AR5d
Last 10 Davs			AR10d

Table 4: Normalized percent prediction error rates for the testbed site pairs for the August 2001 dataset. The figure denotes four categories: (1) prediction based on GridFTP data in isolation (AVG25), (2) regression between GridFTP and NWS network data with the three filling in techniques (G+N), (3) regression between GridFTP and disk I/O data with the three filling in techniques (G+D), and (4) regression based on all three data sources (G+N+D). Shaded portions indicate a "best of class" comparison between the approaches. All percentage values are averages based on different file categories.

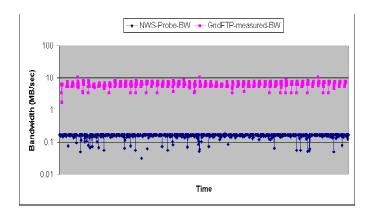
	Only GidFTP Logs [VSF02]	Linear Regression between GridFTP Logs and Network Load [VS02]				ession betwe s and Disk L		Linear Regression uUing aAl Three Data Sources			
	AVG25	G+N NoFill	G+N LV	G+N Avg	G+D NoFill	G+D LV	G+D Avg	G+N+D NoFill	G+N+D LV	G+N+D Avg	
LBL-ANL	24.4%	22.4%	20.6%	20%	25.2%	21.7%	21.4%	22.3%	17.7%	17.5%	
LBL-UFL	15%	18.8%	11.1%	11%	20.1%	11.6%	11.9%	11.1%	8.7%	8%	
ISI-ANL	15%	12%	9.5%	9%	13.1%	13%	11.4%	11%	8.9%	8.3%	
ISI-UFL	21%	21.9%	16%	14.5%	22.7%	19.7%	18.8%	14.7%	13%	12%	
ANL-UFL	20%	21%	20%	16%	21.8%	19.9%	19.3%	15.3%	16.7%	15.5%	

Table 5: Normalized percent prediction error rates for the various site pairs for December 2001 dataset. Figure denotes four categories: (1) prediction based on GridFTP data in isolation (Moving Avg), (2) regression between GridFTP and NWS network data with the three filling in techniques (G+N), (3) regression between GridFTP and disk I/O data with the three filling in techniques (G+D), and (4) regression based on all three data sources (G+N+D). Shaded portions indicate a comparison between our approaches. All percentage values are averages based on different file categories.

	Only GidFTP Logs	O	ession betweend Network		Linear Regro Logs	ession betwe s and Disk L		Linear Regression Using All Three Data Sources		
	Moving Avg	G+N NoFill	G+N LV	G+N Avg	G+D NoFill	G+D LV	G+D Avg	G+N+D NoFill	G+N+D LV	G+N+D Avg
LBL-ANL	20%	23%	17.6%	17%	24%	19.5%	19%	20%	15.2%	15.4%
LBL-UFL	16%	17%	14.7%	13%	16%	14%	14.8%	14.5%	12.2%	12%
ISI-ANL	13%	12%	10.6%	9.8%	12.2%	11.3%	11%	11.3%	9%	8.7%
ISI-UFL	17%	19.3%	13.2%	12%	18%	15%	12%	15%	10%	10.8%
ANL-UFL	18%	18.7%	14.8%	14%	17.8%	17%	16.7%	15.6%	14%	13.3%

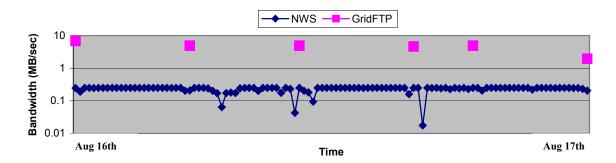


(a) LBL-ANL

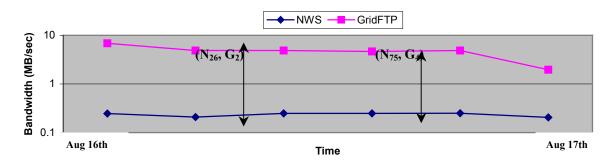


(b) ISI-ANL

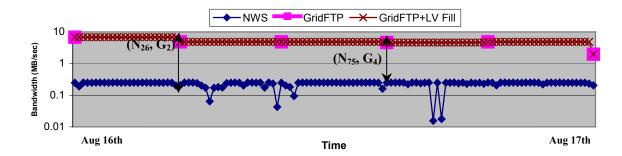
Figure 1: (a) LBL-ANL GridFTP (approximately 400 transfers at irregular intervals) end-to-end bandwidth and NWS (approximately 1,500 probes every five minutes) probe bandwidth for the two-week August'01 dataset. (b) GridFTP transfers and NWS probes between ISI-ANL



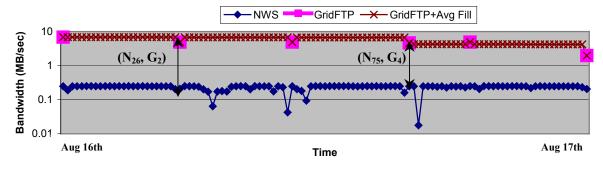
(a) Measured GridFTP and NWS



(b) NoFill



(c) Last Value Filling (LV)



(d) Average Filling (Avg)

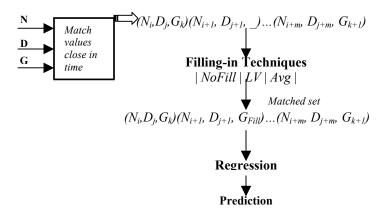


Figure 3: Sequence of events for deriving predictions from GridFTP (G), disk load (D), and NWS (N) datastreams.

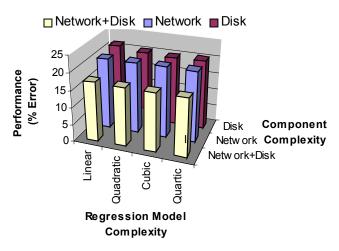


Figure 4: Visualization comparing error, complexity of algorithm, and components included for the site pair LBL and ANL.

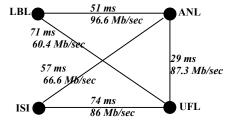


Figure 5: Depiction of network settings for our testbed sites connected through OC-48 network links. For each site pair, roundtrip times and bottleneck bandwidths for the link between them is shown.

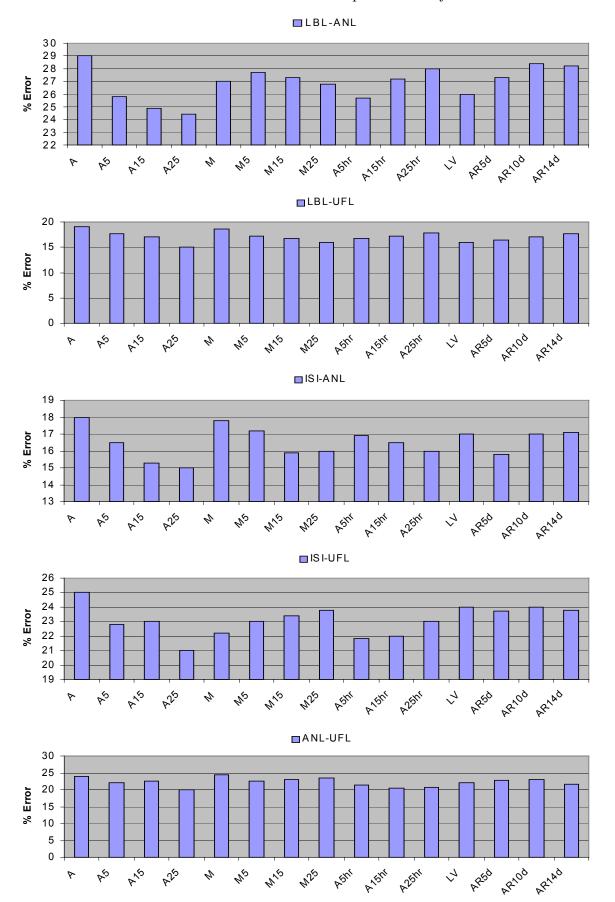


Figure 6: Univariate predictor performance for the testbed site pairs. Predictors include mean-based, median-based, and autoregressive models. The figure also shows context-insensitive variations of all the predictors.

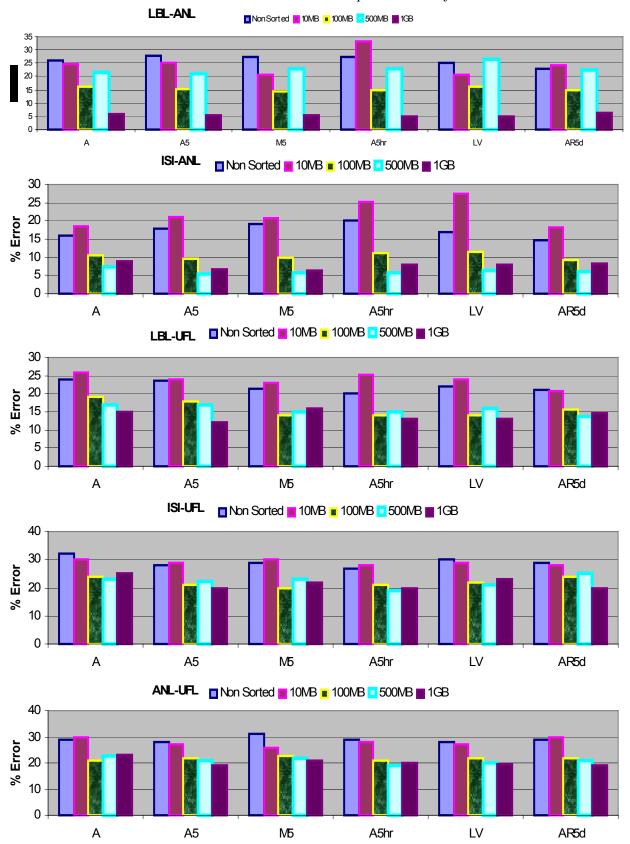


Figure 7: Impact of classification and the reduction in percent error rates for the testbed (context-sensitive filtering).

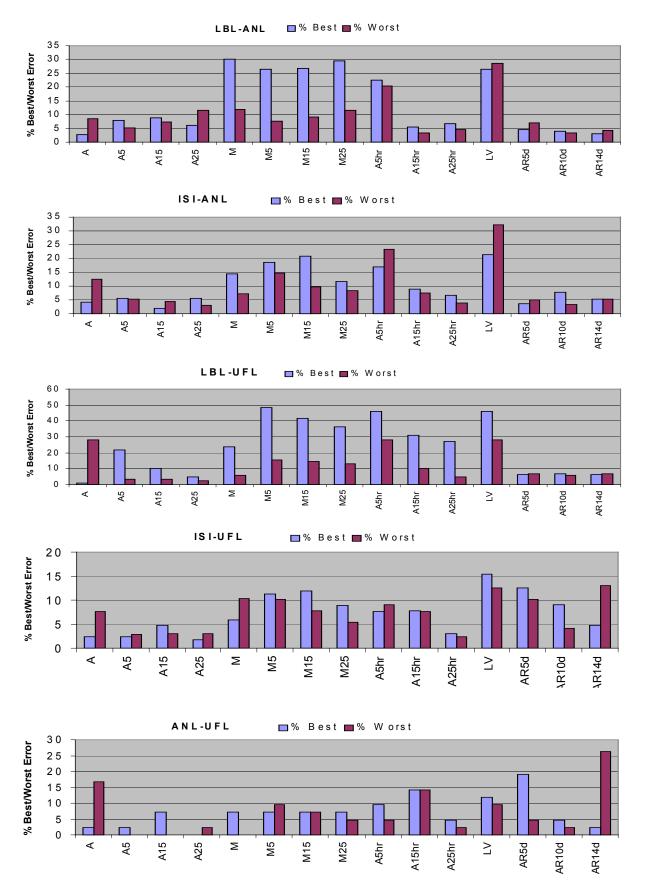
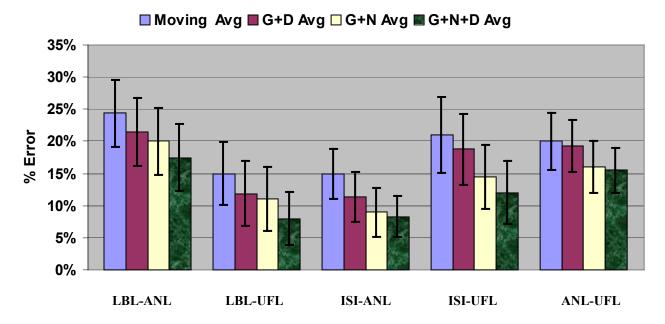


Figure 8: Relative performance of predictors As a percentage best/worst of all predictors for all site pairs.



(a) Comparison of normalized percent errors for the predictors with 95% confidence limits

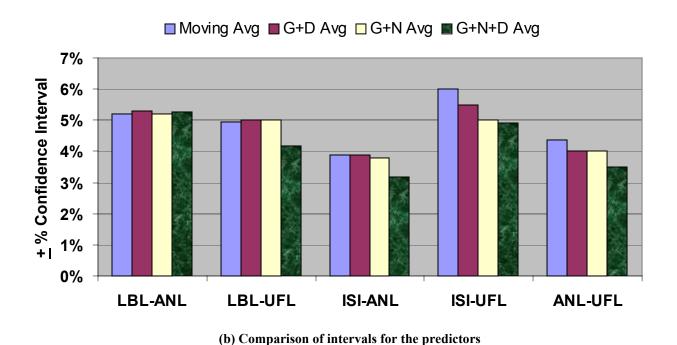


Figure 9: (a) Normalized percent prediction error and 95% confidence limits for August 2001 dataset based on (1) prediction based on GridFTP in isolation (MovingAvg), (2) regression between GridFTP and disk I/O with Avg filling strategy (G+D Avg); (3) regression between GridFTP and NWS network data with Avg filling strategy (G+N Avg), and (4) regressing all three datasets (G+N+D Avg). Confidence Limits denote the upper and lower bounds of prediction error. For instance, the LBL-ANL pair had a prediction range of [17.3% ± 5.2%]. (b) Comparison of the percentage of variability among the predictors.

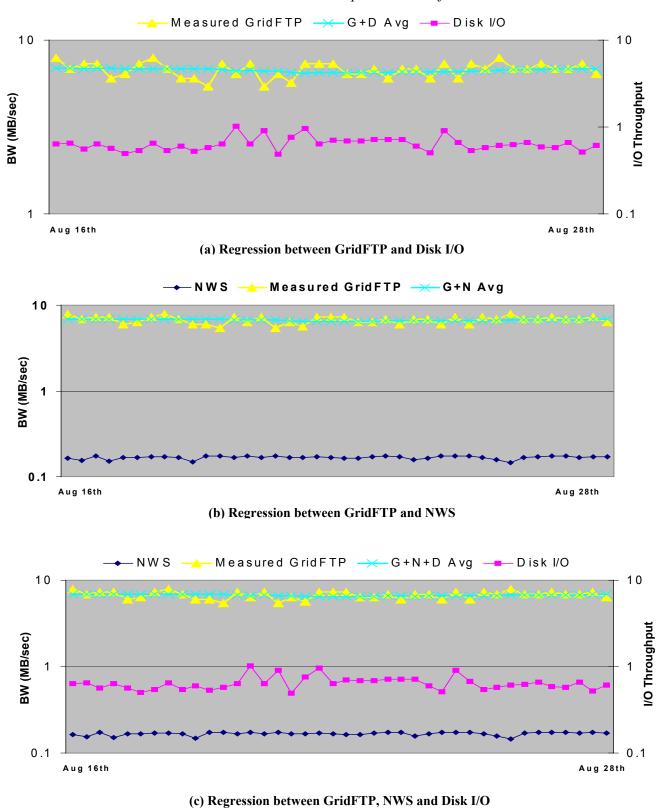


Figure 10: Predictors for 100 MB transfers between ISI and ANL for August 2001 dataset. In the graphs, GridFTP, G+D Avg, G+N+D Avg, and NWS are plotted on the primary *y*-axis; while Disk I/O is plotted on the secondary *y*-axis. I/O throughput denotes transfers per second.

G+D Avg-Linear G+D Avg-Quadratic G+D Avg-Cubic G+D Avg-Quartic

Figure 11: Error rates of polynomial regression models for the G+D Avg predictor for the various site pairs. Polynomials include linear, quadratic, cubic, and quartic models.